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Prakriyāpradarśinī - an open source *subanta* generator

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Abstract: Prakriyāpradarśinī is an attempt to imitate subanta derivation process by prakriyā method given in Siddhāntakaumudī (SK) of Bhaṭṭojī Dikṣita (1910) using an open source PHP code. Our goal is to imitate SK regarding applicability of rules and give the user step by step derivation. The machine handles strīpratyayaprakaraṇa also.

In theory, there is no fixed order of rules for derivation process in sapādasaptādhyāyī of Aṣṭādhyāyī, but if we analyze SK for practical application of rules, rules are applied in some kind of order. The authors have tried to find out the optimum order of application of rules from Sanskrit NLP perspective and are proposing an ‘NLP order model’ and ‘NLP order hypothesis’ for coding subantaprakaraṇa of SK. This is extremely beneficial from coding perspective, because it would decrease the iterations compared to the prevalent ‘conflict resolution model’ e.g. for a 10 step process, in the ‘conflict resolution model’ computer will check whether any of 4000 odd sūtras are applicable or not for 10 times and resolve the conflict i.e. >40000 event checking, whereas in ‘NLP order model’ it would check the criteria for application of sūtras chronologically i.e. only with marginally above 4000 event checking.

The present paper tries to analyze the necessity of user input in subantaprakaraṇa of SK for proper declension. The paper also discusses some of the issues in rule ordering and conflict resolution for Sanskrit NLP from grammatical perspective.

Keywords: Aṣṭādhyāyī, Computational Linguistics, Conflict Resolution, Natural Language Programming, NLP Order Model, NLP Order Hypothesis, Pāṇini, Prakriyā, Prakriyāpradarśinī, Siddhāntakaumudī

1 Introduction

Prakriyāpradarśinī is an attempt to imitate the derivation process given in Siddhāntakaumudī (SK) of Bhaṭṭojī Dīkṣita using an open source PHP code. Our goal is to imitate Siddhāntakumudī in terms of applicability of rules. The main difference between the present approach and earlier approaches for derivation is in the methodology.

The present approach uses ‘NLP order model’ in contrast to ‘Conflict resolution model’ employed earlier. The details about this model will be discussed in section 4. This machine also handles *strīpratyayas*. The other difference is regarding the licence of the code. This code is an open source code which anyone can use and modify according to his specific need. We have spent enough time reinventing the wheel in Sanskrit NLP world. It is high time to move on to an open source world.

2 Review of literature

In all available literature regarding Sanskrit NLP, Goyal et al. (2009) has been found the most relevant to the pursuit at hand, therefore it has been commented upon.

Introduction of that paper mentions that one has to be precise in what one wants to simulate. We have taken SK as base. *Sūtras* and *vārtikas* which are accepted in SK have been incorporated in the code. Section 2 of the same paper raises an issue that Scharf’s (2008) method closely follows SK and not that of Aṣṭādhyāyī (AS). Our present endeavour also is a simulation of SK, but there is kind of a reconciliation of AS method and SK method. The only place where we have some liberty in order is *sapādasaptādhyāyī*¹, whereas order of *tripādī*² is more or less unchangeable for any researcher. We have deployed one do-while loop for most of the rules of *sapādasaptādhyāyī*, which continues till the input and output are same i.e. till there is no rule in *sapādasaptādhyāyī* which can apply now. So, it works like AS method, but if the rules are in random order, the rules will loop over several times (at least 5-6 times as per our estimate). This is a heavy cost on the code and server. It is better if we arrange the rules in the method specified in ‘NLP order model’ which closely follows a *prakriyā* method to economize

¹ AS chapter 1-1 to 8-1

² AS chapter 8-2 to 8-4

70 on time and space. The words ‘NLP order model’ refer to an alternate
 71 order of rules where rules are organized in an order which is more suited for
 72 Computational Linguistics related coding as compared to AS order of rules.

73 While the observation ‘AS models generation’ is true, we would like to
 74 draw attention that sometimes the declension varies according to speaker’s
 75 intention e.g. whether ‘priyatri’ would mean *priyāḥ trayah yasya saḥ* or
 76 *priyāḥ tisrah yasya saḥ* depends on speaker’s intention (SK on 6.4.4). The
 77 derivation also varies according to it. Therefore, it is mandatory that we
 78 take user input on places which we find ambiguous or not amenable to cod-
 79 ing, especially in machines which do single word derivation with no context
 80 whatsoever.

81 According to section 4 of Goyal et al. (2009), simulation of Sanskrit
 82 grammar involves the following factors. (1) Interpretation of sūtras using
 83 the metalanguage described by Pāṇini in the AS, (2) faithful representa-
 84 tion of sūtras, (3) automatic triggering of rules and (4) automatic conflict
 85 resolution. On these parameters, our machine works as mentioned below.

86 1. For **interpretation**, we have used explanation offered by SK.
 87 2. *Sūtras* are **represented** as faithfully as possible. As and when any
 88 wrong output is encountered, the code is re-inspected and necessary correc-
 89 tions are made.

90 3. **Automatic triggering** of rules is done as and when the necessary
 91 conditions are satisfied.

92 4. For **conflict resolution**, *apavāda*, *parasūtra*, *aṅgakārya* and alter-
 93 ation in *pratyayas* are given priority, which takes care of majority of con-
 94 flicts. In case any conflict remains, it also is resolved by altering the order
 95 of the rule application. Unlike Pāṇini’s structure, computer language codes
 96 are almost always executed in a linear sequential manner. Therefore, per-
 97 haps the most favoured mode to stop execution of code is to place the code
 98 in ‘if blocks’. Considering the number of *sūtras*, the places of conflict are
 99 relatively very few, and most of the time, there is ample grammatical liter-
 100 ature to resolve the conflicts. So, in our opinion, finding the correct order
 101 for computer execution is a possibility within human reach. If, after all
 102 possible re-ordering, the conflict still remains, we can add a patch for that
 103 particular word. This is the fundamental theory behind restructuring the
 104 order of rules in ‘NLP order model’.

105 Section 4 of that paper mentions usage of regular expressions to rep-
 106 resent patterns and positions to represent right and left context and alter

string by them. Slight modification in the present system is the use of array rather than string. Its advantage is that the array can store multiple strings e.g. it can store both *vāk*, *vāg* by application of *vā'vasāne* (8.4.56) for future manipulation. If we store the output in a string rather than an array and modify that string by rules of grammar, it becomes difficult to handle optional forms. For rule triggering, we have also used regular expressions and two custom made functions 'sub' and 'arr'.

Section 6 of the same paper highlights that rule *sasajuṣo ruḥ* (8.2.66) is an exception to *pūrvatrāsiddham* (8.2.1) metarule and its implication in coding. SK specifically mentions that *rutva* is not *asiddha* to some rules which require *rutva* as triggering event³. Therefore, we have placed *sasajuṣo ruḥ* (8.2.66) at two places. First place is before the application of these rules. Second is its usual *tripādī* place. To prevent re-application of the *sūtra*, we have remembered that once the rule has been applied. We check while applying for second time whether the rule has been applied already earlier. If yes, we do not apply it again, otherwise the rule in *tripādī* is applied.

Section 7 of the same paper points out a question regarding repeat application of *yāḍāpah* (7.3.113) cyclically in *ramā + ṅe* because of *ṅittva* of *pratyaya*. To circumvent this cyclical application of rules (when the rule should apply only once), in the do-while loop we increase \$start by one, every time the code makes a loop. In the condition for the rule triggering, we ask whether variable \$start is equal to 1 or not. If it is 1, then only the rule applies. Otherwise, it does not apply.

Scharf (2009) has evaluated conflict resolution in AS with four different principles and concluded that 'determining which rule has precedence in the shared domains is not reducible to a single principle'. Therefore, we have adopted the conflict resolution explanations given in SK and coded according to it. Cardona (2009) has analysed the principle of *pūrvatrāsiddham* and its allowable exceptions in grammar. It is mainly in tandem with what traditional grammar texts offer. We have coded properly for it as far as *subanta* generation is concerned.

With this background in mind, let us proceed with the paper.

³ See SK on *ato roraplutādaplute* (6.1.113), SK part 1 page 99.

3 Overview of the project

The project aims at creating an open source PHP code⁴ which would derive noun forms of a given word step by step according to SK. At places, we have also adopted the explanation offered by SK and displayed it to the user to make the derivation easier to understand. But, in some cases, SK explains why a particular rule is not applicable. If it places some constraint on code, we have left out that explanation in display.

The process within the code is in SLP1 transliteration for ease of coding because it assigns a single letter to all *Devanāgarī* characters as has been mentioned by Hyman (2009). This minimizes ambiguity. For example, प्रउग and प्रौढ both would have ‘prau’ in their IAST / HK transliteration, whereas SLP1 for them would be ‘prauga’ and ‘prOQa’. This also eases out transliteration to *Devanāgarī* and other encodings.

User can enter the words in IAST, SLP1 or *Devanāgarī* transliteration. A peculiar problem arises when the user enters a *halanta* word in *Devanāgarī*. Because of variety of input methods, there are sometimes associated white spaces with *halanta* marker. So, before processing, one has to remove the white spaces first with help of code. The *sūtra* display to user is bilingual in Howard Kyoto protocol and *Devanāgarī*. The word under derivation is shown in *Devanāgarī*. At later stage, if need be, this can also be shown in SLP1, IAST or HK or any other transliteration method.

If the machine needs to know some additional parameters based on the word given, ajax.php is called and it gives additional fields in the front end for user to fill. Then ultimately subanta.php is fired and the output is displayed to user.

We have followed the following style pattern, so that the display conveys additional information regarding the *sūtra*: Red colour for headings and error messages, gray colour for *vidhi sūtras* or *apavāda sūtras*, green colour for *paribhāṣās* and *sañjñāvidhāyaka sūtras* and yellow colour for explanatory notes. Thus, the user can get the information regarding *sūtra* type also from the display without much difficulty.

In Sanskrit grammar, there are certain words which are *nityadvivacana* or *nityabahuvacana* or whose derivation is same in all three *lingas*. We have listed some of such words as and when they occur in SK and display the user the information if some *pratyaya* can not apply to this word.

⁴<https://github.com/drdhaval2785/SanskritSubanta>

Certain fonts do not display *upadhmānīya* or *jīhvāmūlīya* properly. Therefore, for proper display we have chosen Siddhanta⁵ font as our default font. We have used ‘!’ to denote anusvāra. A brief note regarding some special characters used in code can be seen here⁶.

Sometimes, we had to do things not mandated by SK to accommodate user tendency. For example, we have observed that the users usually enter *visarga* instead of *sakāra* at the end. We have accepted that user behavior and modified the code to give back the *sakāra* in *prakriyā*.

Discussion in this paper would not be according to the order of *sūtras* in SK or AS, but according to the sequence of code *subanta.php*⁷ in which *sūtras* are applied in this machine.

4 NLP order model and NLP order hypothesis

Though Pāṇini’s rule order is treated as very strict, in our opinion there is a possibility of re-ordering them for ease of computational linguistics. We propose the ‘NLP order model’ and suggest an alternative rule order for *subantaprakaraṇa*⁸.

As the *sūtras* have a kind of grouping based on similarity or conflict, there is some free space in which the *sūtra* order can be moved up / down in machine. Based on our experience, we put forward this ‘**NLP order hypothesis**’ for deciding rule ordering as per ‘**NLP order model**’ for Computational linguistics.

Let us suppose that *sūtras* are in the order

A1, A2, ..., A(k), ..., A(n)

The **Range of freedom** which the *sūtra* A(k) enjoys in term of moving it upwards or downwards in an algorithm for application to input string is equal to the **range (A(min), A(max))**, where **A(min)** is the last previous *sūtra* which can alter the input string for A(k) and make A(k) inapplicable for any possible word in Pāṇini’s grammar. **A(max)** is the first next *sūtra* until which no other *sūtra* can alter the output string of A(k) for

⁵<http://www.svayambhava.org/> www.svayambhava.org/

⁶https://github.com/drhdhaval2785/SanskritSubanta/blob/master/encoding_notes.txt/

⁷<https://github.com/drhdhaval2785/SanskritSubanta/blob/master/subanta.php/>

⁸https://github.com/drhdhaval2785/SanskritSubanta/blob/master/rule_order.txt/

203 any possible word of Pāṇini's grammar. So, it will be possible to decide
 204 the proper position of most of *sūtras* for computational linguistics with this
 205 range in future. If there are still some places where such rule ordering gives
 206 wrong results, patches may have to be applied.

207 Let us clarify how we came to such a conclusion with a working example.

208 If we have a look at the derivation of *rāma* word in SK, the following
 209 *sūtras* are important.

210 *rāma + ŋe -*

211 *svaujasamauṭchaṣṭābhyāmbhisñebhyāmbhyasñasibhyāmbhyasñasosāmñiyossup*

212 (15), *ñeryaḥ* (240) and *supi ca* (291). The numbers in bracket indicate their
 213 position in our code based on 'NLP order model'.

214 Upto this point, we have taught the machine that the order is *svau-*
 215 *jasamauṭchaṣṭābhyāmbhisñebhyāmbhyasñasibhyāmbhyasñasosāmñiyossup* (15),
 216 *ñeryaḥ* (240) and then *supi ca* (291). At this point the range of *sūtra* *ñeryaḥ*
 217 is (15,291).

218 When we move forward in the declension and reach *rāma + bhyas*, there
 219 is a possibility of application of *bahuvacane jhalyet* (290). Now as a coder,
 220 we will have to decide where we should put this rule. At this juncture, we
 221 can see that *supi ca* (7.3.102) is the rule which can alter the output string
 222 (*rāma+bhyas -> rāmā+bhyas*) and render *bahuvacane jhalyet* inapplicable
 223 (no *akāra* at end). Therefore, we can not place *bahuvacane jhalyet* after *supi*
 224 *ca*, otherwise, *bahuvacane jhalyet* will see an altered string (*rāmā+bhyas*)
 225 and it can not apply. Therefore, we have to put it just before *supi ca*
 226 i.e. at place 290. Thus, our code sequence will be *svaujasamauṭchaṣṭāb-*
 227 *hyāmbhisñebhyāmbhyasñasibhyāmbhyasñasosāmñiyossup* (15), *ñeryaḥ* (240),
 228 *bahuvacane jhalyet* (290) and *supi ca* (291). The lower limit upto which
 229 this code can be shifted is *svaujasamauṭchaṣṭābhyāmbhisñebhyāmbhyasñas-*
 230 *sibhyāmbhyasñasosāmñiyossup* (15). Thus we have a range for this *sūtra*
 231 *ñeryaḥ* as (15,290). As the code progresses, the interval gets shortened.

232 Let us take the derivation of 'tad' *pulliṅga*, to see how the range gets
 233 minimized as the code progresses.

234 In this code the relevant rules for our discussion are *tyadādīnāmaḥ*
 235 (7.2.102), *ato guṇe* (6.1.97) and *bahuvacane jhalyet* (7.3.103). When we
 236 want to place *tyadādīnāmaḥ* at its proper place, it has to be before *ato*
 237 *guṇe*, which also should be before *bahuvacane jhalyet*. Therefore their *inter*
 238 *se* applicability would be *tyadādīnāmaḥ -> ato guṇe -> bahuvacane jhalyet*.
 239 Out of which *tyadādīnāmaḥ* and *ato guṇe* are treated in a single piece of

code. Therefore they are applied at specified place let's say 191th place
 looking at other exigencies. As is obvious, they should be placed before
bahuvacane jhalyet, because otherwise the word *tad + bhyas* would not
 have conditions suitable for *bahuvacane jhalyet* to apply (it would have *tad*
 + *bhyas* instead of expected *ta + bhyas*). So, the range of *sūtra neryaḥ*
 has become (191,290). Similarly the range *bahuvacane jhalyet* has become
 (191,291) instead of earlier (15,291). In this way, the location of *sūtra* be-
 comes more and more restricted as the code advances. This way we can
 keep on playing with the location of *sūtras* as SK advances. It is beyond
 our mathematical capacities, but we suggest that if the sequences of appli-
 cation of rules in *prakriyā granthas* like SK are evaluated mathematically,
 near perfect rule order with least possible iteration loops may be derived
 mathematically.

If at any given point, there is a difficulty in identifying proper location
 of a *sūtra*, and alteration in position gives erroneous forms, we can create a
 patch (some code to sort the issue out) for the same and overcome it. But
 in most of the *sūtras* in *subantaprakaraṇa*, it was possible to find their place
 properly⁹.

5 Overview of the code

There are mainly 6 files in our code – *ajax requirements.docx*, *ajax.php*,
function.php, *mystyle.css*, *script.js*, *subanta.html* and *subanta.php*. We shall
 now examine the salient features of the code used in this machine.

5.1 *subanta.html*, *ajax requirements.docx*, *ajax.php* and *script.js*

Because we are treating only the *subantaprakaraṇa* of SK (SK pp. 110-326),
 it is not possible that machine knows all other *sūtras* of AS. Even otherwise,
 there are cases when the word declension depends on user's intention (*vi-*
vakṣā). Therefore, it is not possible that a machine alone can give us desired
 declension without user input. We have tried to make machine responsive
 to the word entered by user. It shows appropriate radio buttons to gather

⁹ Till the paper for *subanta* generation was redrafted, the code has progressed beyond
striṅṅpratyayaprakaraṇa and 80% of *tinantaprakaraṇa*. This 'NLP order model' still held
 good. So, the suggestion made earlier has been substantiated as the code advanced..

प्रकृति: : viSvapA

पुंलिङ्ग स्त्रीलिङ्ग नपुंसकलिङ्ग

Devanagari IAST SLP1

Is this ākārānta dhātu? No

Submit Query

Figure 1

User input window

270 additional information from user, if needed for our purpose. ajax require-
 271 ments.docx, ajax.php and script.js are the scripts responsible for seeking the
 272 input of user.

273 Let us clarify this with an example. If a user enters a word with ‘ā’ at
 274 the end and wants to decline it in masculine gender, we need to get the in-
 275 formation whether this is *ākārānta dhātu* or not. To get this information, we
 276 take user input via ajax when a word ending in ‘ā’ is entered and masculine
 277 gender is selected as shown in Figure 1. For description of methodology,
 278 please see section 11. The detailed list of cases where we take user input are
 279 available on our website¹⁰.

280 At this point let us clarify about two types of user input which have been
 281 deployed. The first type is where the input is taken, because the declension
 282 depends on the intention of user. An example of this case would be whether
 283 ‘sarva’ is used as *sañjñā* or not. This is non-negotiable type of input, in the
 284 sense that even in future we would not be able to do away with them. We

¹⁰[https://github.com/drdhaval2785/SanskritSubanta/blob/master/ajax/
 %20requirement.docx/](https://github.com/drdhaval2785/SanskritSubanta/blob/master/ajax/%20requirement.docx/)

285 have written ‘no’ in the list in this file¹¹ for this type of feedback. The second
 286 type is where we do not know (as of now) how to decide some parameter
 287 e.g. right now we are asking the user whether the word is *ābanta* or not. In
 288 future, when the machine learns how to differentiate a word having *ābanta*,
 289 we will no longer need this input. In that sense, this type is negotiable. As
 290 the machine progresses, these input points can be removed. We have noted
 291 down the second type of input with a ‘yes’ and a note on how we can remove
 292 them in future. Veracity or otherwise of this list is open to suggestion.

293 5.2 function.php

294 Pāṇinian grammar works on many *sūtras* which are called for execution as
 295 and when a condition is satisfied. We have devised some functions based on
 296 those *sūtras* for repetitive work e.g. functions ‘prat’ (*pratyāhāra*), savarna,
 297 vridhhi, guna, dirgha, ṭi, mit etc. This file¹² also holds the data sets e.g.
 298 vowels, consonants etc in addition to the functions.

299 5.3 slp-dev.php and dev-slp.php

300 They are transliterator codes which convert SLP1 transliteration to *Devanā-*
 301 *garī* and vice versa. These codes are borrowed and modified from Dicrunch
 302 code of Ananda Loponen¹³.

303 5.4 subanta.php

304 This is the code which actually processes the entered word and shows the
 305 result back to the user. The most intricate part about automatic declension
 306 machine have always been rule triggering, conflict resolution and ordering
 307 of *sūtras*. The first two issues will be dealt with at their respective places.
 308 The third is a bit lengthy, so its concept is discussed in the section of rule
 309 ordering. Details of rule order are available on our website¹⁴.

¹¹https://github.com/drdhaval2785/SanskritSubanta/blob/master/user_input.pdf

¹²<https://github.com/drdhaval2785/SanskritSubanta/blob/master/function.php/>

¹³<https://github.com/drdhaval2785/sanskrit/tree/master/diCrunch/>

¹⁴https://github.com/drdhaval2785/SanskritSubanta/blob/master/rule_order.txt/

310 This code was first developed as a *sandhi* generator, and later on merged
 311 with *subanta* generator. Therefore, coding for all *sūtras* mentioned in *sand-*
 312 *hiprakaraṇa* is kept intact within this *subanta* generation machine. Most
 313 of them are even used for derivation processes also. User may note some
 314 *praṅghya* related portion which is not relevant to *subanta* generation directly
 315 in the code, but we have retained it from the legacy of *sandhi* generator.

316 For most of the code we have retained a ‘+’ sign between *prakṛti* and
 317 *pratyaya*. But when it comes to *dvitvaprakaraṇa* in *tripādī*, ‘+’ sign creates
 318 some disturbance in the function because of a coding issue which we could
 319 not overcome. So, for now we have removed + sign before *dvitvaprakaraṇa*.
 320 There are not many *sūtras* after *dvitvaprakaraṇa*, which need the identifica-
 321 tion of *pada* and *pratyaya*. Therefore, there is not much information which
 322 is lost by removing this ‘+’ sign.

323 *khari ca* (8.4.55) is a cyclic process till all possible combinations are
 324 over. So, we have kept a while loop till there is no member satisfying the
 325 condition. Let us clarify it with an example.

326 e.g. *suhṛd + sup* => *suhṛdd + su* (by *anaci ca* (8.4.47)) => *suhṛdt + su*
 327 (by *khari ca* (8.4.55)).

328 At this stage, there still is a ‘d’ preceding ‘t’. Therefore, *khari ca* (8.4.55)
 329 finds its application once again. So, we have kept a while loop till there is
 330 no member satisfying the condition. By doing this, we could ensure that
 331 ‘d’-> ‘t’ transition can still take place, and gives *suhṛttsu*.

332 There is some gray area regarding application of *paribhāṣā parjanyaaval-*
 333 *lakṣaṇapravṛttiḥ* (*pa* 119)¹⁵. If there is a combination of *car* + *khar* letter,
 334 should *khari ca* (8.4.55) apply, because ‘*car*’ is itself a subset of ‘*jhal*’?
 335 Though there is no difference in the form, the rule must apply because of
 336 *parjanyaavallakṣaṇapravṛttiḥ* (*pa* 119) *paribhāṣā*, because the conditions for
 337 application of rule are satisfied. Should we display such cases or not is yet
 338 to be determined, but, anyway, the code is mature enough to handle both
 339 the choices.

340 6 Variables

341 AS uses variables very effectively in its structure. Several *sañjñās* are as-
 342 signed to the word and they are made use of at a later stage e.g. ‘*sarva*’ gets

¹⁵ The *paribhāṣā* numbers refer to *paribhāṣenduśekhara* (1913)

343 *sarvanāma-sañjñā* by *sarvādīni sarvanāmāni* (1.1.27) and AS uses them at
 344 places like *sarvanāmnaḥ smai* (7.1.14). In coding parlance, its close approx-
 345 imation is something like this:

```
346 if ( input word = sarva ) { $sarvanama=1; }
```

347 This assigns \$sarvanama value of 1 like Pāṇini assigns it *sarvanāma*
 348 *sañjñā*.

```
349 if ($sarvanama===1) { Do sarvanāmnaḥ smai (7.1.14) }.
```

350 This checks whether \$sarvanama is equal to 1 or not, and executes the
 351 code. It is similar to application of *sarvanāmnaḥ smai* (7.1.14) in case the
 352 word has *sarvanāma sañjñā* in Pāṇinian system.

353 Thus, variables play very crucial role in the simulation of Pāṇini's gram-
 354 mar. We have enumerated some of the variables which we have used in our
 355 code so that the reader may have a bird's view about what is happening in
 356 the code.

357 Examples of variables used in the code are: *sup*, *pada*, *bha*, input word,
 358 gender, transliteration, *nadī*, *nī*, *ābanta*, *taddhita*, *dhātu* etc. All the vari-
 359 ables can be seen in function.php and subanta.php. Their explanation and
 360 importance are given in the code itself as and when they are applied for the
 361 first time. Unless specified otherwise, the meaning of different values are: 0
 362 - no application, 1 - mandatory application and 2 - optional application.

363 Variables are used for two purpose in our code: (1) to remember that
 364 some rule **has been applied** e.g. \$Ap=1 would mean that the word is
 365 derived from some *āp pratyaya* (that rule has already applied) and (2) to
 366 remember that some rule **has to be applied** e.g. \$sarvafinal=1 would mean
 367 that all rules specific to *sarvanāmas* have to be applied. We have chosen
 368 the variable names close to the corresponding grammatical notation so that
 369 the it is easy to understand the code.

370 7 Rule Triggering

371 There are specific *prakriyās* to be followed in grammar when specific con-
 372 ditions are satisfied. Therefore, appropriate rule triggering is of utmost
 373 importance for success of the code. We have used inbuilt functions of PHP,
 374 syntax of PHP, operators and some user defined functions to check whether
 375 the conditions for application of a rule are met or not.

8 Sample Code

We will show a sample code here along with its explanation so that the user may get the feel of what is happening in the background for fetching the required output.

8.1 Rule triggering code

```

381 /* ñeryaḥ (7.1.13) */
382 if (arr($text,'/[a][+][ñ][e]/') && $pada=== "pratyaya" && $so ===
383 "ñe" )
384 {
385 $text = one(array("a+ñe"),array("a+ya"),0);
386 echo "<p class = sa >By ñeryaḥ (7.1.13) :</p>";
387 echo "<p class = sa > ( . . ) :</p>";
388 display(3);
389 $ñe=1; // 0 - This sūtra has not applied. 1 - This sūtra has been applied.
390 }

```

8.2 Rule triggering explanation

1. if (arr(\$text,'/[a][+][ñ][e]/') && \$pada=== "pratyaya" && \$so === "ñe")

In this section we check the following conditions – the suffix (\$so) is 'ñe', *akāra* is followed by *ñe* and *ñe* is a *pratyaya*. When these conditions are satisfied, the rest of the code is executed.

2. { }

The bracketed area is code which is to be executed.

3. \$text = one(array("a+ñe"),array("a+ya"),0);

In this section, we convert 'a+ñe' to 'a+ya' i.e. we apply *ñeryaḥ sūtra*.

4. echo "<p class = sa >By *ñeryaḥ* (7.1.13) :</p>";

echo "<p class = sa > (. .) :</p>";

We display the *sūtra* which has been applied in this case.

5. display(3);

We display the word to user. (In *Devanāgarī*)

6. \$ñe=1;

We remember that the *ñeryaḥ* has been applied to this word, for future use in code.

9 Conflict Resolution

We have used usual Pāṇinian methods like *apavāda*, *parasūtra* etc. for conflict resolutions as discussed in SK. Grammar books and their commentaries provide plenty of literature on conflict resolution. There are many *paribhāṣās* also. The most important *paribhāṣā* dealing with conflict resolution is *paranītyāntaraṅgāpavādānāṃ uttarottaraṃ balīyaḥ* (pa 38).

This has been taken care of implicitly in rule ordering itself. We have tried to place *parasūtra*, *nītya prakriyās*, *antaraṅga prakriyās* and *apavādasūtras* before *pūrvasūtra*, *anītya prakriyās*, *bahiraṅga prakriyās* and *utsargasūtras* respectively. If there is conflict among *para*, *nītya*, *antaraṅga*, *apavāda*, the later wins. Such encounter has not happened in *subanta* generation stage. Sometimes *paribhāṣās* are *nītya* / *anītya*. Sometimes *pūrvavipratīṣedha*¹⁶ applies e.g. *numaciratṛjvadbhāvebhyo nuṭ pūrvavipratīṣedhena* (vā 4374). We have used *pūrvavipratīṣedha* whenever it is explicitly mentioned in the text of SK. Thus, generic application of metarules is not possible. So they will be applied in specific cases only. In addition, there are places where grammarians have difference of opinion. In such cases we have taken SK as authority¹⁷. Whatever has been accepted in SK is imitated in the code. If SK is silent on some topic, other available commentaries are explored for solution.

9.1 Methods to avoid application of a rule

1. Ordering *apavāda*, *parasūtra*, *antaraṅga*, *nītya*, *aṅgavidhi*, *pratyaya* alteration rules before the *utsarga*, *pūrvasūtra*, *bahiraṅga*, *anītya* and other rules. This way, the later group sees a modified string which does not satisfy criteria for their application.
2. Remembering that a rule has to be applied in future e.g. variable \$purvapara=1 means that the rule *pūrvaparāvaradaksiṅottarāparādharāṇi vyavasthāyāmasaṃjñāyām* (1.1.33) will apply later on. When the turn of this rule comes we check whether variable \$purvapara is equal to 1 or not.

¹⁶ *pūrvavipratīṣedha* means processes where the preceding *sūtra* is given priority over subsequent *sūtra*, violating the general rule.

¹⁷ See SK part 1 page 161 under the rule *trestrayaḥ* 7.1.53

439 3. Remembering that a rule is not to be applied in future e.g. if *ato'm*
 440 (7.1.24) has been applied, we store the value of variable \$atom as 1
 441 and when conditions of application of *svamornapumṣakāt* (7.1.23) are
 442 tested, we tell it not to apply the *sūtra* to words where *ato'm* (7.1.24)
 443 has been applied.

444 9.2 Notes on issues in conflict resolution

445 With this background, let us examine some of the issues which cropped up
 446 during process of simulating *subantaprakaraṇa*:

447 1. *ārambhasāmarthya*.

448 *supi ca* (7.3.102) is *parasūtra* compared to *nāmi* (6.4.3). Even then, it
 449 does not apply in case of *nāmi* (6.4.3) even though it is *parasūtra* be-
 450 cause of *ārambhasāmarthya*. So, due care needs to be taken in coding
 451 such cases according to the explanation given in grammar texts.

452 2. *anityatva* of *paribhāṣās*.

453 For example, *sannipātaparibhāṣā* does not apply in case of *ne pratyaya*.
 454 *supi ca* (7.3.102) applies in that case. *kaṣṭāya kramaṇe* (3.1.14) is an
 455 example of *anityatva*¹⁸ of this *paribhāṣā*. If we code for such *paribh-*
 456 *hāṣās* to apply in every case, this form will get distorted. The better
 457 alternative is to apply such *paribhāṣās* only in cases where SK has
 458 validated its applicability.

459 3. *sasajuṣo ruḥ* and treatment of *sakāra*.

460 Following code blocks should be placed before *sasajuṣo ruḥ* (8.2.66)
 461 for execution of code, otherwise their ultimate *sakāra* may take
 462 *sasajuṣo ruḥ* (8.2.66) to give undesired forms. - 1. *etattadoḥ*
 463 *sulopo'ko'nañsamāse halī* (6.1.132), 2. *so'ci lope cetpādapūraṇam*
 464 (6.1.134), 3. *aniditām hala upadhāyāḥ kṛiti* (6.4.24) for *sraṃs*, *dvaṃs*
 465 etc 4. *vasusraṃsudhvaṃsvanaḍuhām daḥ* (8.2.72). These code blocks
 466 have to be kept before actual execution of *sasajuṣo ruḥ* (8.2.66).

467 4. *ṇatva*.

¹⁸ This means the same rule applies with some condition at some place, but does not apply at some other place even if the same condition is satisfied.

It is difficult to identify where *pada* ends and another starts, especially in *samāśas*. Therefore, the conflict resolution among *sūtras ekājuttarapade ṇaḥ* (8.4.12) / *raṣābhyāṃ no ṇaḥ samānapade* (8.4.1) / *aṭkuvāṇnumvyavāye'pi* (8.4.2) is extremely difficult until *samāsa* is parsed effectively.

10 Notes on some grammar issues and their coding implications.

This section is devoted to brief mention of issues we encountered during development of this software, which may serve as a reference point to the future researchers in case they face the same difficulty in implementation.

1. Do-while loop for *sapādasaptādhyāyī*.

pūrvatrāsiddham (8.2.1) *sūtra* creates two separate data spaces for *sūtras* of AS namely (1) *sapādasaptādhyāyī* and (2) *tripādī*. There are two other *sūtras* which are also used to create separation of data spaces like *asiddhavad atrābhāt* (6.4.22) and *ṣatvatukorasiddhaḥ* (6.1.86)¹⁹. We have not treated these two *sūtras* here, because they are not explicitly treated in *subantaprakaraṇa* of SK. So we will keep our discussion focused on *pūrvatrāsiddham* (8.2.1) only²⁰. To put it in very basic terms, *pūrvatrāsiddham* (8.2.1) makes provision that the rules in *sapādasaptādhyāyī* have no fixed order of application and those of *tripādī* have to be applied after all possible *sapādasaptādhyāyī* rules have applied to the word. Even inside *tripādī*, the rules are to be applied sequentially. For *sapādasaptādhyāyī*, we have created a do-while loop which checks whether the word entering and word coming out of this loop is the same or not. It continues till both are the same (i.e. until there is no rule in *sapādasaptādhyāyī* which can apply and alter the word). In most of the cases, no looping of *sapādasaptādhyāyī* is needed. In other cases, only one looping of *sapādasaptādhyāyī* is needed. This is quite low burden on code as compared to checking 4000 rules every time. After that, it moves on to *tripādī*, where the

¹⁹ See Goyal et al. (2009) section 3.6

²⁰ ‘na lopaH supsvarasajjhātugvidhiṣu’ is dealt with in the code at places where it is applicable.

498 rules are usually written in sequence of their appearance in AS, so that
499 the rules apply sequentially.

500 In this approach, there is a limitation. There are certain rules in
501 *sapādasaptādhyāyī*, which can apply only once. For such rules, we
502 have applied them before do-while loop or we have specified in do-while
503 loop that this rule is to be applied when the code is being executed
504 for the first time only by specifying \$start=1. When the code comes
505 for execution for second time, we check whether \$start==1. In the
506 second loop \$start has value of 2 and therefore, this rule is not applied
507 for second time. This is, in a nutshell, what we do to ensure secrecy
508 of data space of *sapādasaptādhyāyī* and *tripādī*.

509 2. *asiddhatva* of *ṣakāra* for applicability of *sasajuṣo ruḥ* (8.2.66).

510 There are two possible cases (1) *dhātus* - *pipaṭhiṣ*, *āśiṣ*. (2) not *dhātus*
511 - *doṣ*, *dhanuṣ*. The usual user tendency is to enter *ṣakāra* at the end
512 of such words. If we proceed with this word, there is no *sakāra* at the
513 end of the word, and therefore *sasajuṣo ruḥ* (8.2.66) can not apply. To
514 circumvent this, *ṣakāra* has to be converted back to *sakāra* for *sasajuṣo*
515 *ruḥ* (8.2.66)'s application. Similar situation appears in case of *vivakṣ*,
516 *didhakṣ*, *pipakṣ* etc. *skoḥ samyogādyorante ca* (8.2.29) does not apply
517 because *kutva* is *asiddha* to *kalopa*. Similarly, in case of *cikīrṣ* - *ṣatva*
518 is *asiddha* to *rātsasya* (8.2.24). Therefore, *rātsasya* (8.2.24) sees *cikīrṣ*
519 only and not *cikīrṣ*. Thus, it causes elision of last *ṣakāra*. For such
520 typical cases, patches have to be applied in the code.

521 3. All *pratyaya* alterations have to be completed before applying *sūtras*
522 which remove *it* marker, otherwise it would not be possible to check the
523 exact *pratyaya* e.g. *rāma+ne* -> *rāma+ya* by *ñeryaḥ* (7.1.13) *sūtra*.
524 If we had removed the *it* marker 'ñ' before application of *ñeryaḥ*, the
525 situation would have been *rāma+e* and we would have tough time
526 finding out whether this is a *ñit pratyaya* or not.

527 4. attributes of stem and suffix.

528 It is of vital importance to remember the attributes of stem and suffix
529 e.g. *sarvanāma*, *sarvanāmasthāna*, *it* markers, previous application
530 of some *sūtra*, *bhāṣitapumskatva*, *trjvadbhāva*, original input word,
531 *sambuddhi*, *ṣaṭ*, *nadī sañjñā* etc. Some of them are discussed here in

brief. These attributes may be absent, present or optionally present in grammar. Because of these three types, it is not advisable to use boolean variables for attributes, as they cater to only presence and absence.

Status of *sarvanāma* - It is important to know whether the word is a *sarvanāma*, not a *sarvanāma* or optionally a *sarvanāma*. With the help of *sarvādi gaṇa* and some user input, we decide whether \$sarvafinal = 0 (not *sarvanāma*), 1 (*sarvanāma*) or 2 (optionally *sarvanāma*). Sometimes, we need to add a member at a later stage in a *gaṇa*. e.g. ‘*sva*’ is pronoun only if it is used not in the meaning of *jñāti* or *dhana*. If we add ‘*sva*’ directly into *sarvanāma* set, it will give erroneous results if it is used in sense of *jñāti* or *dhana*.

it markers play an important role in derivation process e.g. *aco ṅṅiti* (7.2.115) will apply only when the *pratyaya* has ‘*ṅ*’ or ‘*ṅ*’ as *it* marker. Similarly there are many rules where we have to know about *it* marker.

It is wise to remember the first input word. This may be needed after some time e.g. *hali lopaḥ* (7.2.113) mandates elision of ‘*id*’ of ‘*idam*’ when certain conditions are met. Till this stage, ‘*idam*’ is already converted to ‘*ida*’ by *tyadādīnāmah* (7.2.102) and *ato guṇe* (6.1.96). Therefore we need to check whether this ‘*ida*’ was derived from ‘*idam*’ or not.

sambuddhi forms are different than regular forms. Therefore, it is mandatory to remember whether the *pratyaya* is *sambuddhi* or not. e.g. *eṅhrasvātsambuddheḥ* (6.1.69) will only apply in case of *sambuddhi*.

5. *stoḥ ścunā ścuḥ* (8.4.40) and *śāt* (8.4.44).

This case is different form of expression than regular ABC-> ADC context based transformation. Therefore, special treatment is needed. In this case, no specific order of letters is mandated. This rule applies in case of juxtaposition rather than order. Therefore, a separate code is needed to handle this rule. Similarly *ṣṭunā ṣṭuḥ* (8.4.41) and *na padāntāṭṭoraṇām* (8.4.42), *anāmnavatīnagarīṅāmīti vācyam* (*vā* 5016) and *toḥ ṣiḥ* (8.4.43) need specific treatment.

6. *na lumatāṅgasya* (1.1.63).

566 This rule prevents conversion of *padānta kim* to *ka*. Explanation of
 567 SK is given a place in derivation scheme. *na lumatāṅgasya* (1.1.63)
 568 also bars application of *pratyayalope pratyayalakṣaṇam* (1.1.62) - so
 569 we have to remember whether *luk* has happened or not. We should
 570 also remember that *na lumatāṅgasya* (1.1.63) is *anitya paribhāṣā*. It
 571 does not apply in case of *tricituroḥ striyāṃ tisṛcatasṛ* (7.2.99)²¹.

572 7. *sthānivadbhāva*.

573 There is a great deal of literature on what is *sthānivadbhāva* and what
 574 is not in SK. Therefore, we have not treated *sthānivadbhāva* generi-
 575 cally. We have coded according to it only when the derivation demands
 576 such intervention to be made.

577 8. *nimittāpāye naimittikasyāpāyaḥ*.

578 When doing elision by *saṃyogāntasya lopaḥ* (8.2.23), we have to be
 579 ready for application of *nimittāpāye naimittikasyāpāyaḥ*.

580 9. Aberrant behaviour of rules.

581 *vṛddhyauttvatṛjvadbhāvaguṇebhyo num pūrvavipratīṣedhena* (vā 4373)
 582 mandates that the *numāgama* gets precedence over rules mentioned
 583 here by *pūrvavipratīṣedha*. Similarly, *numaciratṛjvadbhāvebhyo nuṭ*
 584 *pūrvavipratīṣedhena* (vā 4374) mandates that *nuḍāgama* gets prece-
 585 dence over rules mentioned here by *pūrvavipratīṣedha*. These rules are
 586 exception to the general precedence of *parasūtra*.

587 Sometimes grammarians try to justify the derivation of a *śiṣṭa* word by
 588 ‘*akṛtavīyūhāḥ pāṇinīyāḥ*’. An example of it can be seen in application
 589 of *acaḥ* (6.4.138).

590 10. Difference of opinion amongst grammarians. There is a difference of
 591 opinion among grammarians regarding *prarīṇām*²² and *varṣābhū*. We
 592 have accepted SK’s position.

593 11. Difficulty in coding.

594 *rutva* can happen any time in between code. Therefore,
 595 *upadeśe’janunāsika it* (1.3.2) which elides *ukāra* of ‘*su!*’ might not

²¹ See pp. 196 of SK part 1.

²² See SK part 1 p. 220

work properly. Therefore, a special patch is made for *rutva* to convert it to *repha*.

ḍhralope pūrvasya dīrgho'ṇaḥ (6.3.111). It is mandatory to remember the *repha* / *ḍhakāra* where this is to be applied. Otherwise, all the *hrasva* + *r/ḍh* combinations in the word would be converted to *dīrgha* + *r/ḍh*. For this purpose, we have added an artificial sign # before *repha* or *ḍhakāra* where this rule has to be applied.

bhāṣītapuṃskatva is very difficult to know from words, and many *prakriyās* depend on whether the word is *bhāṣītapuṃska* or not. This calls for user input.

In case of verbs, users may adopt different conventions for writing a verb. For example, user may insert *ancu*, *añcu*, *anc*, *añc*, *ancu!*, *añcu!* anything. Such behavior is seen more frequently in case of verbs with anubandhas and especially with verbs having a nasal letter in it. So, the code has to be resilient enough to account for such aberrant behavior of users.

Keeping *ato guṇe* (6.1.97) applicable to each *sūtra* creates many issues like interfering with *akāḥ savarṇe dīrghaḥ* (6.1.101). Right now, *ato guṇe* (6.1.97) is used in close conjunction with the rules where grammar textbook mandates it and not as a separate code block.

11 User input

By the words **user input**, we mean ‘getting desired input from user for correct declension of a word’. For a good code, this has to be at bare minimum to enhance user experience. So we have decided what user input fields are negotiable ones and which are not in this file²³. For user input, we use radio buttons. In future, if there is a case where we must check multiple parameters simultaneously, multiple check boxes can also be employed. The documentation for user input is stored in *ajax requirements.docx* e.g. if *ajax.php* uses `$_GET['cond1_2']===2`, it means that condition 1.2.2 in the document is satisfied. The user can easily make out from code and *ajax requirements.docx* what we check out in user input.

²³https://github.com/drdhaval2785/SanskritSubanta/blob/master/user_input.pdf

627

12 Limitations

628 After discussion on methodology followed by us, it would be of interest if we
 629 place before the researchers some problem areas which we encountered, so
 630 that they may be explored further with grammar texts and solutions may
 631 be arrived at.

632 1. Difficulty in identifying attributes. Correct derivation depends on cor-
 633 rect identification of attributes, but it becomes extremely difficult to
 634 identify those attributes in some cases. In such cases, user input may
 635 be our only hope.

636 *nityastrīṅgatva* has to be taken as user input, because it is difficult to
 637 guess *nityastrīṅgatva* by merely looking at the word. Some detailed
 638 analysis of feminine words may help in this regard.

639 *iyaṅnuvaṅsthānatva*²⁴ presupposes the knowledge whether the word
 640 meets the criteria for application of ‘*yaṅ*’ or ‘*vaṅ*’. It is difficult
 641 to analyze this beforehand.

642 There is a perpetual problem whether (1) one should retain “*ni*”, “*ne*”
 643 etc till the *prakriyās* with *nittva* start or (2) should we convert them
 644 to ‘*i*’, ‘*e*’ and remember that it has ‘*n*’ as it. Right now, the former
 645 method is used preferably. *upasarjanībhūtatva* / *pradhānatva* – It is
 646 very important to know these qualities in *samāsas*. Right now, we are
 647 not able to analyse *samāsas*, so we have taken user input in this case.

648 As *kvīn* / *kvīp* *pratyayas* do not leave any mark on the word, they
 649 are difficult to identify. Right now, more or less we are listing out
 650 *kvīnnanta* and *kvībanta* words manually, which may not be workable
 651 in long run. We will have to think of some alternative to overcome
 652 this problem.

653 It is difficult to differentiate *ābantatva* or *ākārānta dhātu* from merely
 654 looking at the word. We will have to understand *ābanta pratyayas* as
 655 well as *dhātu prakriyās* to code properly for them.

656 *abhyastatva*. It is difficult to code for *abhyastatva*, till *abhyāsa prakriyā*
 657 is taught.

²⁴ AS 1.4.4

2. Problems with *dhātus*. It is difficult to identify *dhātus*. Even if we en-
list all *dhātus* in *dhātupāṭha*, there are *nāmadhātus* and *sanādi dhātus*
too, which make it difficult to identify where *kṛdatiṅ* (3.1.93) is to be
applied. Right now, we are looking for ‘*vy*’ to identify *nāmadhātus*.
e.g. *eḍakīyati*. It may need further revision when *nāmadhātus* are
taught to the system.

One needs to find all possible *dhātus* starting with “*ṛ*” to decide
whether *upasargāḍṛti dhātau* (6.1.91) is applicable or not. Even then,
special treatment for *nāmadhātus* is needed because the rule is optional
for *nāmadhātus*.

It is equally difficult to identify *prātipadikas*, because there are many
pratyayas which may derive a new noun from a *dhātu*. Some of them do
not even leave a mark morphologically like *kvīp* / *kvīn* etc. Therefore,
it is really difficult to identify *prātipadikas* and separate them from
dhātus.

3. Issues of morphologic similarity.

Sometimes the *prakriyā* is specified for word ending with some word
e.g. if a *prakriyā* for ‘*ahan*’ is specified and we search for string ‘*ahan*’
only, the *prakriyā* for ‘*sūryāhan*’ may not work well. Therefore, we
have to think about any morphological change which the word might
undergo under the influence of rules of *sandhi* too.

As we work with string of letters in coding, it is difficult to isolate
words ending with ‘*han*’ for applicability of *inhanpūṣāryamṇām śau*
(6.4.12). ‘*han*’ at the end of a word can also be part of ‘*ahan*’, where
this rule would be erroneously applicable. User input or exhaustive
enumeration will be needed for clarity.

4. Issues of *ekādeśa*.

ādyantavadekasmin (1.1.21) rule is difficult to code, because right now
we are keeping a ‘*+*’ sign in between the stem and suffix. It is difficult
in current scheme of things to code properly to remember that the
ekādeśa behaves as end of the previous one and the start of the next
one. Another question which deserves attention is how should *ekādeśa*
be displayed? e.g. *ādguṇaḥ* (6.1.87) mandates *ekādeśa*. What should
we display in case of ‘*rāma+i*’? ‘*rām+e*’ or ‘*rāme+*’? Let us show

- 692 our approach with example. In case of *ādguṇaḥ* (6.1.87), the term
 693 ‘āt’ means that the *ādeśa* is after *akāra*. So, we have kept it ‘rām+e’,
 694 whereas *akāḥ savarṇe dīrghaḥ* (6.1.101) mandates replacement of ‘ak’.
 695 So, we have kept ‘rāmā+t’.
- 696 5. Issues in contextual derivation. As we are not working with sentences
 697 for now, it is difficult to analyse attributes which depend on sentences
 698 e.g. whether there is *anvādeśa* or not in case of derivation of *asmad*
 699 / *yuṣmad*. Currently only words are being treated and not sentences,
 700 so *pāda* related functions are not applied for now.
- 701 6. Different derivation in different meanings. e.g. the word *sudhī* can
 702 be analysed as *suṣṭhū dhīryasyāḥ*, *suṣṭhu dhyāyati*, *suṣṭhu dhīḥ*. The
 703 derivation differs in all the situations. Therefore it is mandatory to
 704 take user input to specify which of these meanings he intends to use.
- 705 7. Listings. Various lists (over and above *gaṇas*) are needed for proper
 706 declension of a word e.g. *ugit dhātus*, *idit dhātus*, *ṛkārānta* words etc.
 707 Exhaustive lists remain to be made for such words.

708 13 Scope for Future work

- 709 1. We have left out accents e.g. *caturanaḍuhorāmudāttaḥ* (7.1.98) – we
 710 have coded only for ‘ām’ and left out ‘udāttaḥ’. We will have to treat
 711 the accent at a later stage for sure, because the *striṛpratyayas* and
 712 *taddhitapratyayas* have very peculiar effect on accent, otherwise mor-
 713 phologically *ṛīp*, *ṛīṣ*, *ṛīn* give the same forms.
- 714 2. The *sūtras* which we have not coded for are specifically mentioned in
 715 code *subanta.php*. The user is advised to refer to them for further
 716 details.
- 717 3. *sūtras* which involve interpretation of *samāsas* are right now on user
 718 input mode. Once *samāsa* interpretation is taught to the machine,
 719 they can be properly coded.

- 720 4. We have prepared a list with a hint whether the requirement for user
 721 input can be done away with or not²⁵. This can serve as a guide for
 722 future researchers. The future attempts should be primarily focused
 723 on removing the unnecessary user input from the machine. Once this is
 724 achieved, researchers can take up the more challenging task of handling
 725 the non negotiable kind of user input.

726 14 Conclusion

727 AS has a complex system of interrelated rules. Various authors have tried in
 728 past to reorganize the order of AS for *prakriyāgranthas*. Similarly there is a
 729 need to reorder the AS for easy implementation of computational simulation
 730 of AS. NLP order model and NLP order hypothesis presented in the present
 731 work is a step in that direction.

²⁵https://github.com/drdhaval2785/SanskritSubanta/blob/master/user_input.pdf

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754 **Appendix : Sample derivation of ramā**

755 Attached below is the derivation of ‘ramā’ word ²⁶

²⁶as on 27.9.2014 from http://lanover.com/lan/sanskrit/subanta.php?first=ram&gender=f&tran=Devanagari&cond2_1=2&step=2_1_2

You entered: रमा + सुँ [Go Back](#)

arthavadadhAturapratyayaH prAtipadikam (1.2.45), kRttaddhitasamAsAzca (1.2.46), pratyayaH (3.1.1), parazca (3.1.2), GyAppradipadikAt (4.1.1), svaujasamauTCaSTAbhyAmbhisGebhyAmbhyasGasibhyAmbhyasGasosAmGy (4.1.2), vibhaktizca (1.4.104) and supaH (1.4.103) :

अर्थवदधातुरप्रत्ययः प्रातिपदिकम् (१.२.४५), कृत्तद्धितसमासाश्च (१.२.४६), प्रत्ययः (३.१.१), परश्च (३.१.२) इयाप्रातिपदिकात् (४.१.१), स्वौजसमौढ्वाभ्याम्भिस्ङेभ्याम्भ्यस्ङसिभ्याम्भ्यस्ङसोसाम्ङ्योस्सुप् (४.१.२), विभक्तिश्च (१.४.१०४) तथा सुपः (१.४.१०३) :

1 - रमा+सुँ

dvyekayordvivacanaikavacane (1.4.22) :

द्व्येकयोर्द्विवचनैकवचने (१.४.२२) :

1 - रमा+सुँ

By suDanapuMsakasya (1.1.43) :

सुडनपुंसकस्य (१.१.४३) :

1 - रमा+सुँ

By upadeze'janunAsika it (1.3.2) :

उपदेशेऽजनुनासिक इत् (१.३.२) :

1 - रमा+सुँ

By tasya lopaH (1.3.9) :

तस्य लोपः (१.३.९) :

1 - रमा+स्

By na vibhaktau tusmAH (1.3.4) :

न विभक्तौ तुस्माः (१.३.४) :

1 - रमा+स्

Figure 2

Sample Derivation: Part 1 of 2

By apRkta ekAlpratyayaH (1.2.41) :

अपृक्त एकाल्प्रत्ययः (१.२.४१) :

1 - रमा+स्

By halGyAbbhyo dIrgHAtsutisyapRktaM hal (6.1.68) :

हल्ङ्याब्भ्यो दीर्घात्सुतिस्यपृक्त हल् (६.१.६८) :

1 - रमा+

By aNo'pragRhyasyAnunAsikaH (8.4.57) :

अणोऽप्रगृह्यस्यानुनासिकः (८.३.५७) :

1 - रमा

2 - रमाँ

Final forms are :

आखिरी रूप हैं :

1 - रमा

2 - रमाँ

Figure 3

Sample Derivation: Part 2 of 2